

INTRODUCTION TO COMPILER DIRECTIVES WITH OPENACC

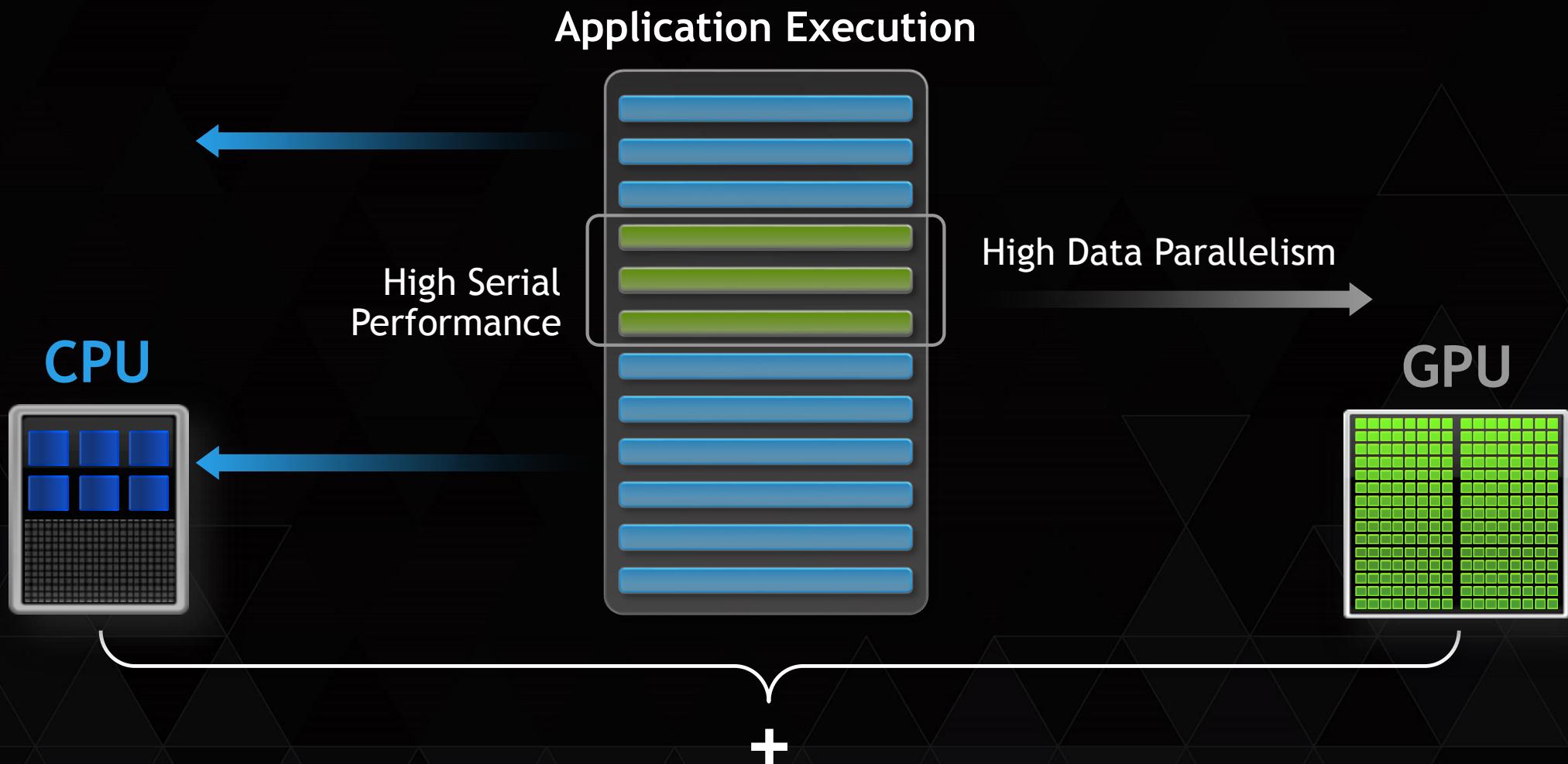
JEFF LARKIN, NVIDIA DEVELOPER TECHNOLOGIES

AGENDA

- ▶ Fundamentals of Heterogeneous & GPU Computing
- ▶ What are Compiler Directives?
- ▶ Accelerating Applications with OpenACC
 - ▶ Identifying Available Parallelism
 - ▶ Exposing Parallelism
 - ▶ Optimizing Data Locality
- ▶ Misc. Tips
- ▶ Next Steps

HETEROGENEOUS COMPUTING BASICS

WHAT IS HETEROGENEOUS COMPUTING?



LOW LATENCY OR HIGH THROUGHPUT?



LATENCY VS. THROUGHPUT

F-22 Raptor

- 1500 mph
- Knoxville to San Jose in 1:25
- Seats 1



Boeing 737

- 485 mph
- Knoxville to San Jose in 4:20
- Seats 200



LATENCY VS. THROUGHPUT

F-22 Raptor

- Latency – 1:25
- Throughput – $1 / 1.42 \text{ hours} = 0.7 \text{ people/hr.}$



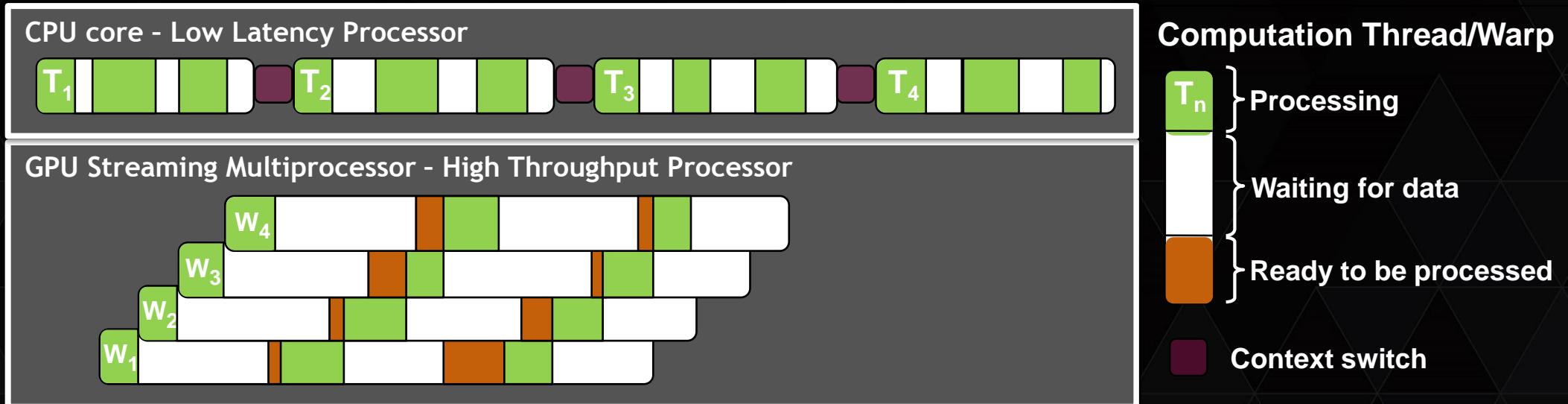
Boeing 737

- Latency – 4:20
- Throughput – $200 / 4.33 \text{ hours} = 46.2 \text{ people/hr.}$



LOW LATENCY OR HIGH THROUGHPUT?

- ▶ CPU architecture must minimize latency within each thread
- ▶ GPU architecture hides latency with computation from other threads



ACCELERATOR FUNDAMENTALS

- ▶ We must expose enough parallelism to fill the device
 - ▶ Accelerator threads are slower than CPU threads
 - ▶ Accelerators have orders of magnitude more threads
 - ▶ Accelerators tolerate resource latencies by cheaply context switching threads
- ▶ Fine-grained parallelism is good
 - ▶ Generates a significant amount of parallelism to fill hardware resources
- ▶ Coarse-grained parallelism is bad
 - ▶ Lots of legacy apps have only exposed coarse grain parallelism

3 APPROACHES TO HETEROGENEOUS PROGRAMMING

Applications

Libraries

Compiler
Directives

Programming
Languages

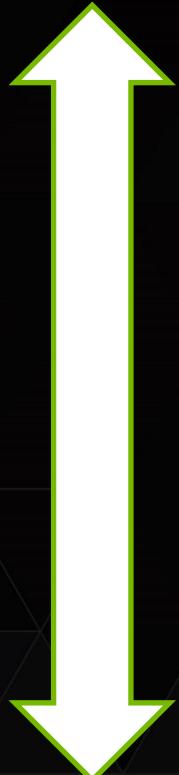
Easy to use
Most Performance

Easy to use
Portable code

Most Performance
Most Flexibility

SIMPLICITY & PERFORMANCE

Simplicity



Performance

- ▶ Accelerated Libraries
 - ▶ Little or no code change for standard libraries, high performance.
 - ▶ Limited by what libraries are available
- ▶ Compiler Directives
 - ▶ Based on existing programming languages, so they are simple and familiar.
 - ▶ Performance may not be optimal because directives often do not expose low level architectural details
- ▶ Parallel Programming languages
 - ▶ Expose low-level details for maximum performance
 - ▶ Often more difficult to learn and more time consuming to implement.

WHAT ARE COMPILER DIRECTIVES?

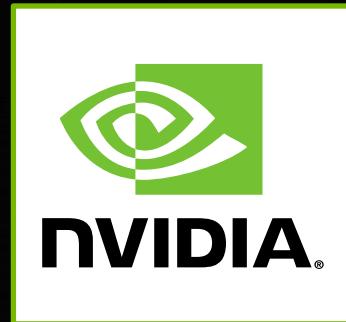
WHAT ARE COMPILER DIRECTIVES?

```
int main() {  
  
    do_serial_stuff()  
  
    for(int i=0; i < BIGN; i++)  
    {  
        ...compute intensive work  
    }  
  
    do_more_serial_stuff();  
  
}
```

Programmer inserts compiler hints.

Execution Begins on the CPU.

Compiler Generates GPU Code.



Data and Execution returns to the CPU.

OPENACC: THE STANDARD FOR GPU DIRECTIVES

- ▶ **Simple:** Directives are the easy path to accelerate compute intensive applications
- ▶ **Open:** OpenACC is an open GPU directives standard, making GPU programming straightforward and portable across parallel and multi-core processors
- ▶ **Portable:** GPU Directives represent parallelism at a high level, allowing portability to a wide range of architectures with the same code.



OPENACC MEMBERS AND PARTNERS



TOTAL



TECHNISCHE
UNIVERSITÄT
DRESDEN



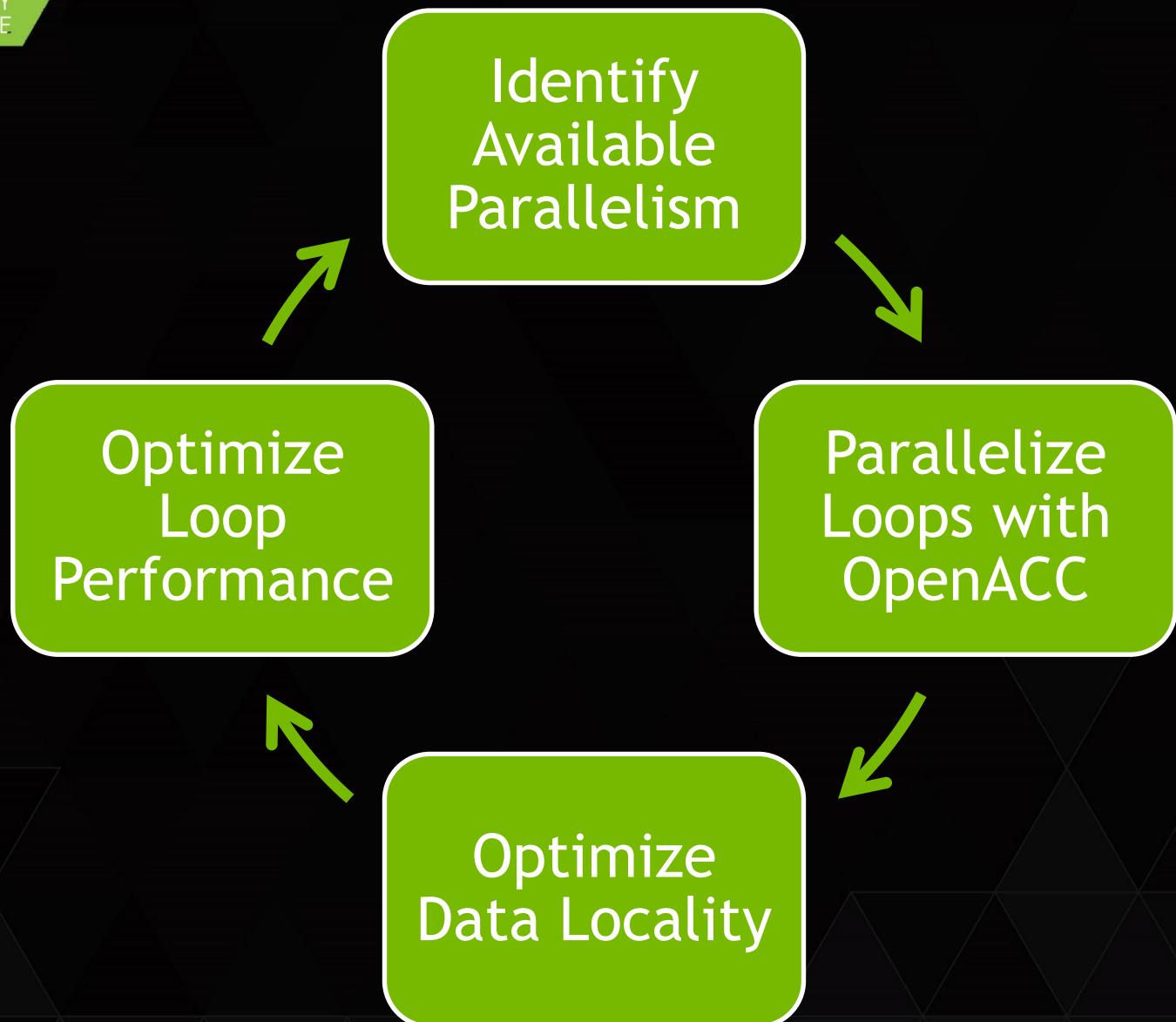
CSCS



LOUISIANA STATE UNIVERSITY

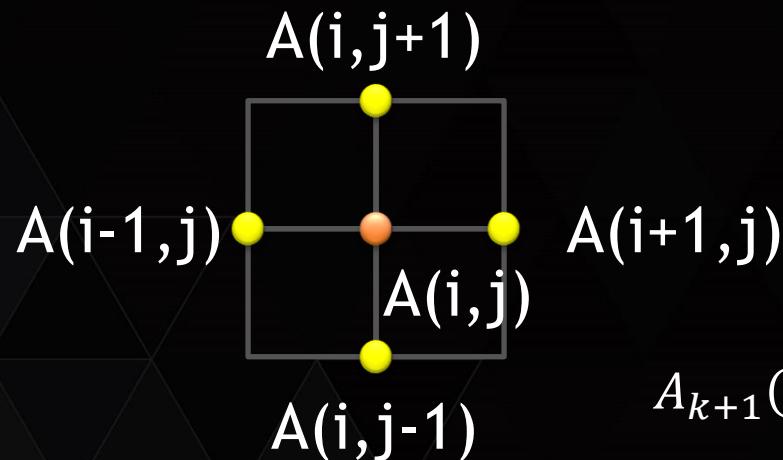


ACCELERATING APPLICATIONS WITH OPENACC



EXAMPLE: JACOBI ITERATION

- ▶ Iteratively converges to correct value (e.g. Temperature), by computing new values at each point from the average of neighboring points.
 - ▶ Common, useful algorithm
 - ▶ Example: Solve Laplace equation in 2D: $\nabla^2 f(x, y) = 0$



$$A_{k+1}(i, j) = \frac{A_k(i - 1, j) + A_k(i + 1, j) + A_k(i, j - 1) + A_k(i, j + 1)}{4}$$

JACOBI ITERATION: C CODE

```
while ( err > tol && iter < iter_max ) {  
    err=0.0;  
  
    for( int j = 1; j < n-1; j++) {  
        for(int i = 1; i < m-1; i++) {  
  
            Anew[j][i] = 0.25 * (A[j][i+1] + A[j][i-1] +  
                                  A[j-1][i] + A[j+1][i]);  
  
            err = max(err, abs(Anew[j][i] - A[j][i]));  
        }  
    }  
  
    for( int j = 1; j < n-1; j++) {  
        for( int i = 1; i < m-1; i++ ) {  
            A[j][i] = Anew[j][i];  
        }  
    }  
  
    iter++;  
}
```



Iterate until converged



Iterate across matrix elements



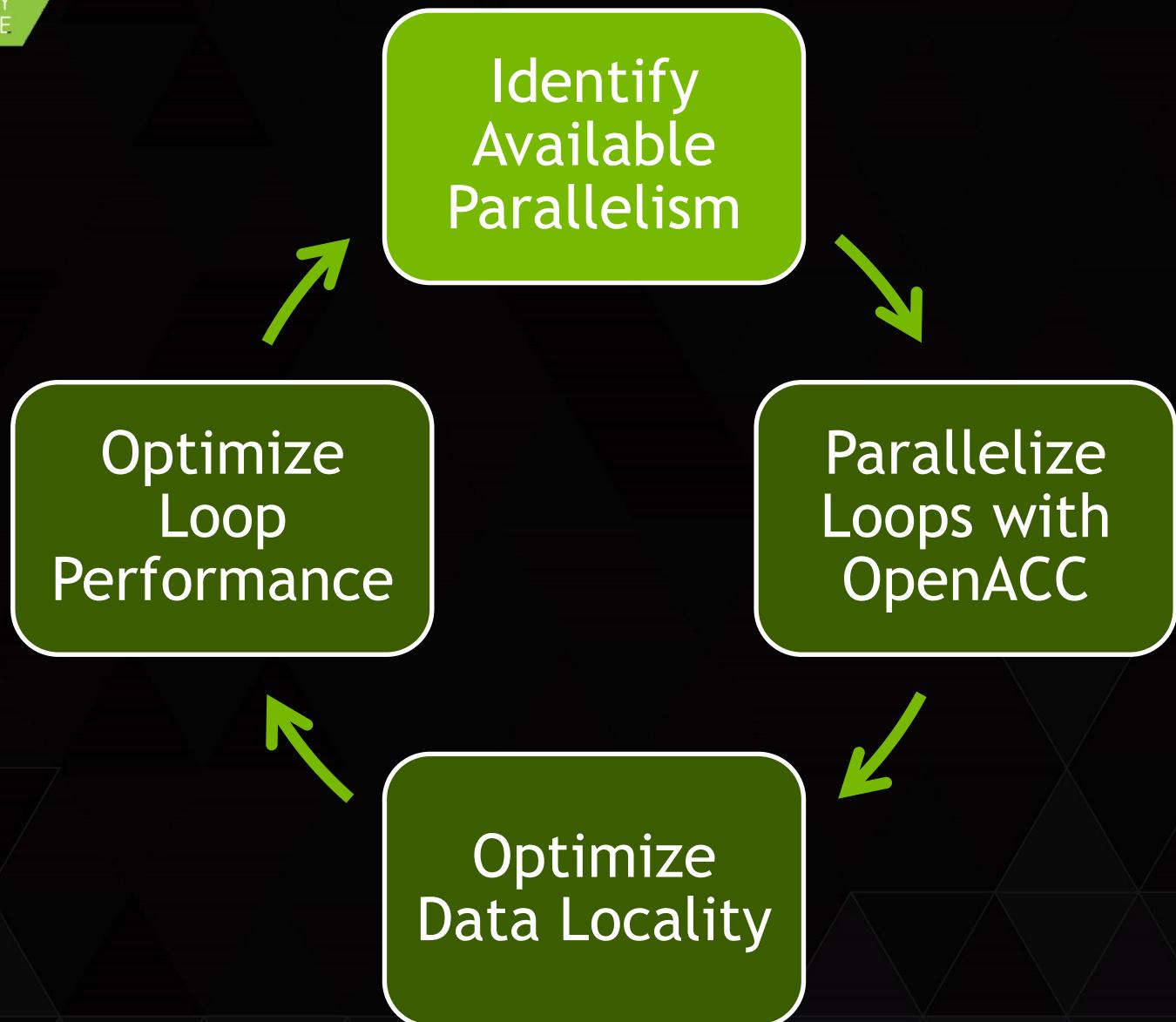
Calculate new value from neighbors



Compute max error for convergence



Swap input/output arrays

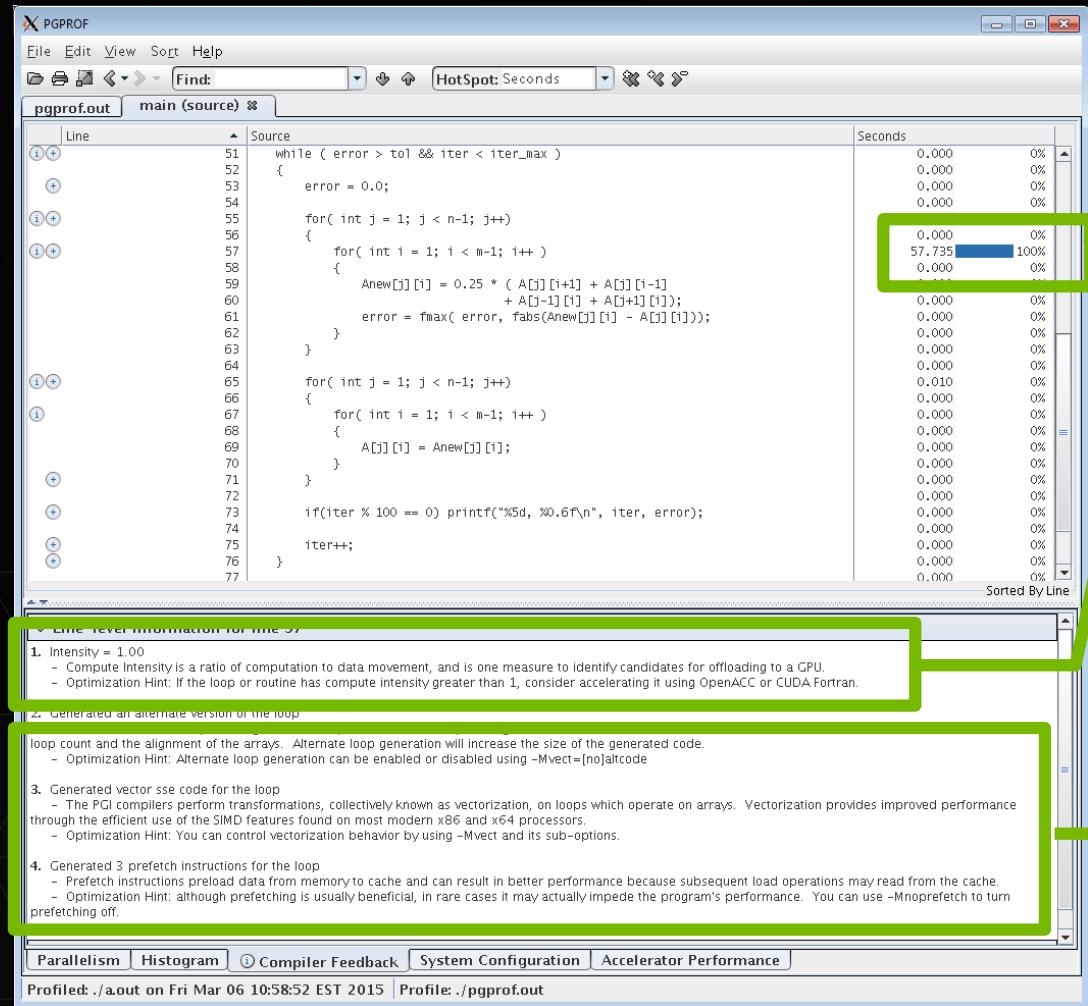


IDENTIFY AVAILABLE PARALLELISM

- ▶ A variety of profiling tools are available:
 - ▶ gprof, pgprof, Vampir, Score-p, HPCToolkit, CrayPAT, ...
- ▶ Using the tool of your choice, obtain an application profile to identify hotspots
- ▶ Since we're using PGI, I'll use pgprof

```
$ pgcc -fast -Minfo=all -Mprof=ccff laplace2d.c
main:
    40, Loop not fused: function call before adjacent loop
        Generated vector sse code for the loop
    57, Generated an alternate version of the loop
        Generated vector sse code for the loop
        Generated 3 prefetch instructions for the loop
    67, Memory copy idiom, loop replaced by call to __c_mcopy8
$ pgcollect ./a.out
$ pgprof -exe ./a.out
```

IDENTIFY PARALLELISM WITH PGPROF



PGPROF informs us:

1. A significant amount of time is spent in the loops at line 56/57.
2. The computational intensity (Calculations/Loads&Stores) is high enough to warrant OpenACC or CUDA.
3. How the code is currently optimized.

NOTE: the compiler recognized the swapping loop as data movement and replaced it with a memcpy, but we know it's expensive too.

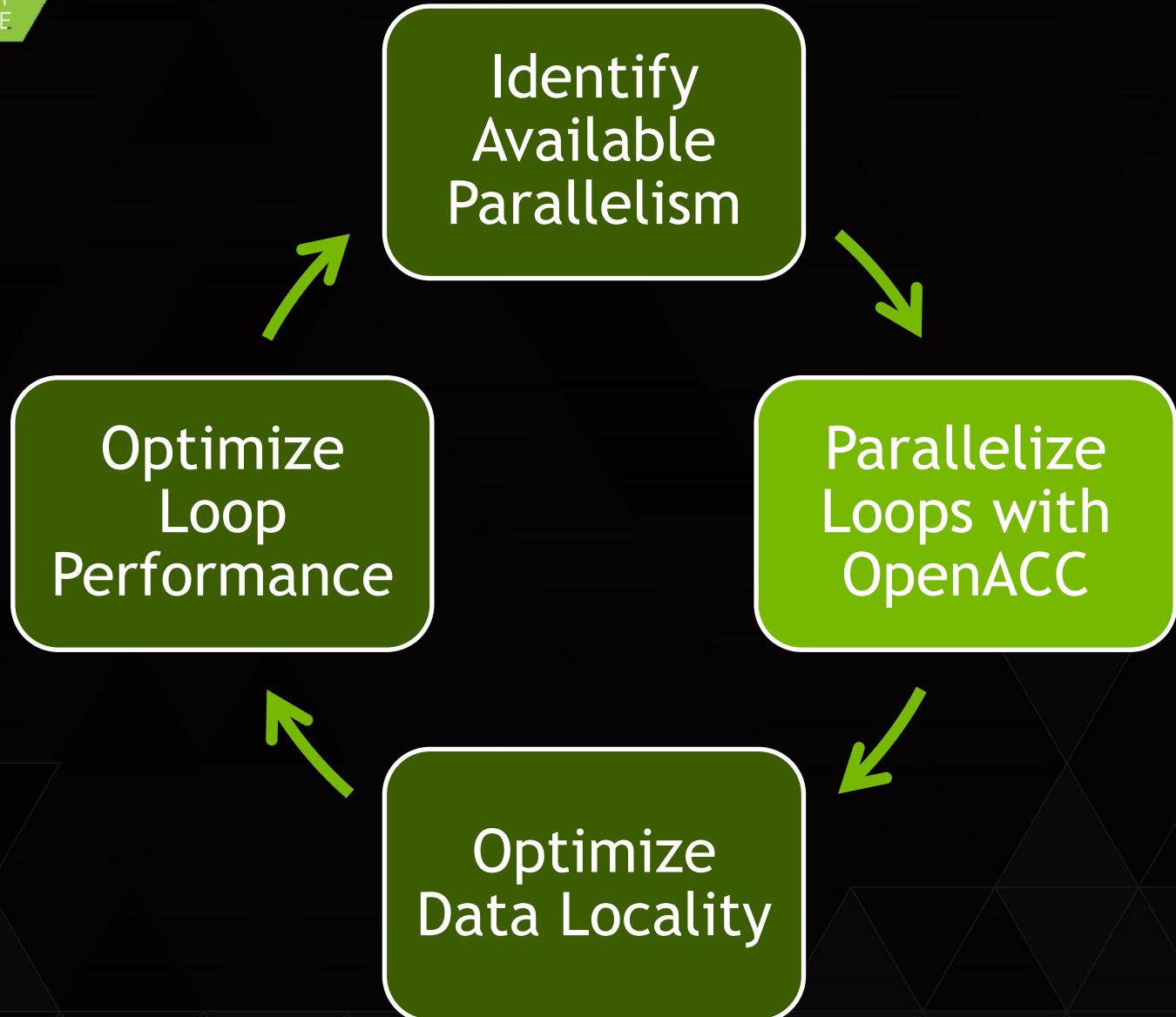
IDENTIFY PARALLELISM

```
while ( err > tol && iter < iter_max ) {  
    err=0.0;  
  
    for( int j = 1; j < n-1; j++) {  
        for(int i = 1; i < m-1; i++) {  
  
            Anew[j][i] = 0.25 * (A[j][i+1] + A[j][i-1] +  
                                  A[j-1][i] + A[j+1][i]);  
  
            err = max(err, abs(Anew[j][i] - A[j][i]));  
        }  
    }  
  
    for( int j = 1; j < n-1; j++) {  
        for( int i = 1; i < m-1; i++ ) {  
            A[j][i] = Anew[j][i];  
        }  
    }  
  
    iter++;  
}
```

Data dependency
between iterations.

Independent loop
iterations

Independent loop
iterations



OPENACC DIRECTIVE SYNTAX

- ▶ C/C++

```
#pragma acc directive [clause [,] clause] ...
```

...often followed by a structured code block

- ▶ Fortran

```
!$acc directive [clause [,] clause] ...
```

...often paired with a matching end directive surrounding a structured code block:

```
!$acc end directive
```



Don't forget acc

OPENACC PARALLEL LOOP DIRECTIVE

parallel - Programmer identifies a block of code containing parallelism. Compiler generates a **kernel**.

loop - Programmer identifies a loop that can be parallelized within the kernel.

NOTE: parallel & loop are often placed together

```
#pragma acc parallel loop
for(int i=0; i<N; i++)
{
    y[i] = a*x[i]+y[i];
}
```

Parallel
kernel

Kernel:
A function that runs
in parallel on the
GPU

PARALLELIZE WITH OPENACC

```
while ( err > tol && iter < iter_max ) {  
    err=0.0;  
  
#pragma acc parallel loop reduction(max:err)  
    for( int j = 1; j < n-1; j++) {  
        for(int i = 1; i < m-1; i++) {  
  
            Anew[j][i] = 0.25 * (A[j][i+1] + A[j][i-1] +  
                                  A[j-1][i] + A[j+1][i]);  
  
            err = max(err, abs(Anew[j][i] - A[j][i]));  
        }  
    }  
  
#pragma acc parallel loop  
    for( int j = 1; j < n-1; j++) {  
        for( int i = 1; i < m-1; i++ ) {  
            A[j][i] = Anew[j][i];  
        }  
    }  
  
    iter++;  
}
```

Parallelize loop on accelerator

Parallelize loop on accelerator

* A *reduction* means that all of the N*M values for err will be reduced to just one, the max.

OPENACC LOOP DIRECTIVE: PRIVATE & REDUCTION

- ▶ The **private** and **reduction** clauses are not optimization clauses, they may be required for correctness.
- ▶ **private** – A copy of the variable is made for each loop iteration
- ▶ **reduction** – A reduction is performed on the listed variables.
 - ▶ Supports +, *, max, min, and various logical operations

BUILDING THE CODE

```
$ pgcc -fast -acc -ta=tesla -Minfo=all laplace2d.c
main:
 40, Loop not fused: function call before adjacent loop
    Generated vector sse code for the loop
 51, Loop not vectorized/parallelized: potential early exits
 55, Accelerator kernel generated
    55, Max reduction generated for error
    56, #pragma acc loop gang /* blockIdx.x */
    58, #pragma acc loop vector(256) /* threadIdx.x */
 55, Generating copyout(Anew[1:4094][1:4094])
    Generating copyin(A[:, :])
    Generating Tesla code
 58, Loop is parallelizable
 66, Accelerator kernel generated
    67, #pragma acc loop gang /* blockIdx.x */
    69, #pragma acc loop vector(256) /* threadIdx.x */
 66, Generating copyin(Anew[1:4094][1:4094])
    Generating copyout(A[1:4094][1:4094])
    Generating Tesla code
 69, Loop is parallelizable
```

OPENACC KERNELS DIRECTIVE

The kernels construct expresses that a region *may contain parallelism* and *the compiler determines* what can safely be parallelized.

```
#pragma acc kernels
{
    for(int i=0; i<N; i++)
    {
        x[i] = 1.0;
        y[i] = 2.0;
    }

    for(int i=0; i<N; i++)
    {
        y[i] = a*x[i] + y[i];
    }
}
```



The compiler identifies
2 parallel loops and
generates 2 kernels.

PARALLELIZE WITH OPENACC KERNELS

```
while ( err > tol && iter < iter_max ) {  
    err=0.0;  
  
#pragma acc kernels  
{  
    for( int j = 1; j < n-1; j++) {  
        for(int i = 1; i < m-1; i++) {  
  
            Anew[j][i] = 0.25 * (A[j][i+1] + A[j][i-1] +  
                                  A[j-1][i] + A[j+1][i]);  
  
            err = max(err, abs(Anew[j][i] - A[j][i]));  
        }  
    }  
  
    for( int j = 1; j < n-1; j++) {  
        for( int i = 1; i < m-1; i++ ) {  
            A[j][i] = Anew[j][i];  
        }  
    }  
}  
  
iter++;  
}
```



Look for parallelism
within this region.

BUILDING THE CODE

```
$ pgcc -fast -acc -ta=tesla -Minfo=all laplace2d.c
main:
 40, Loop not fused: function call before adjacent loop
    Generated vector sse code for the loop
 51, Loop not vectorized/parallelized: potential early exits
 55, Generating copyout(Anew[1:4094][1:4094])
    Generating copyin(A[:, :])
    Generating copyout(A[1:4094][1:4094])
    Generating Tesla code
 57, Loop is parallelizable
 59, Loop is parallelizable
    Accelerator kernel generated
    57, #pragma acc loop gang /* blockIdx.y */
    59, #pragma acc loop gang, vector(128) /* blockIdx.x threadIdx.x */
    63, Max reduction generated for error
 67, Loop is parallelizable
 69, Loop is parallelizable
    Accelerator kernel generated
    67, #pragma acc loop gang /* blockIdx.y */
    69, #pragma acc loop gang, vector(128) /* blockIdx.x threadIdx.x */
```

OPENACC PARALLEL LOOP VS. KERNELS

PARALLEL LOOP

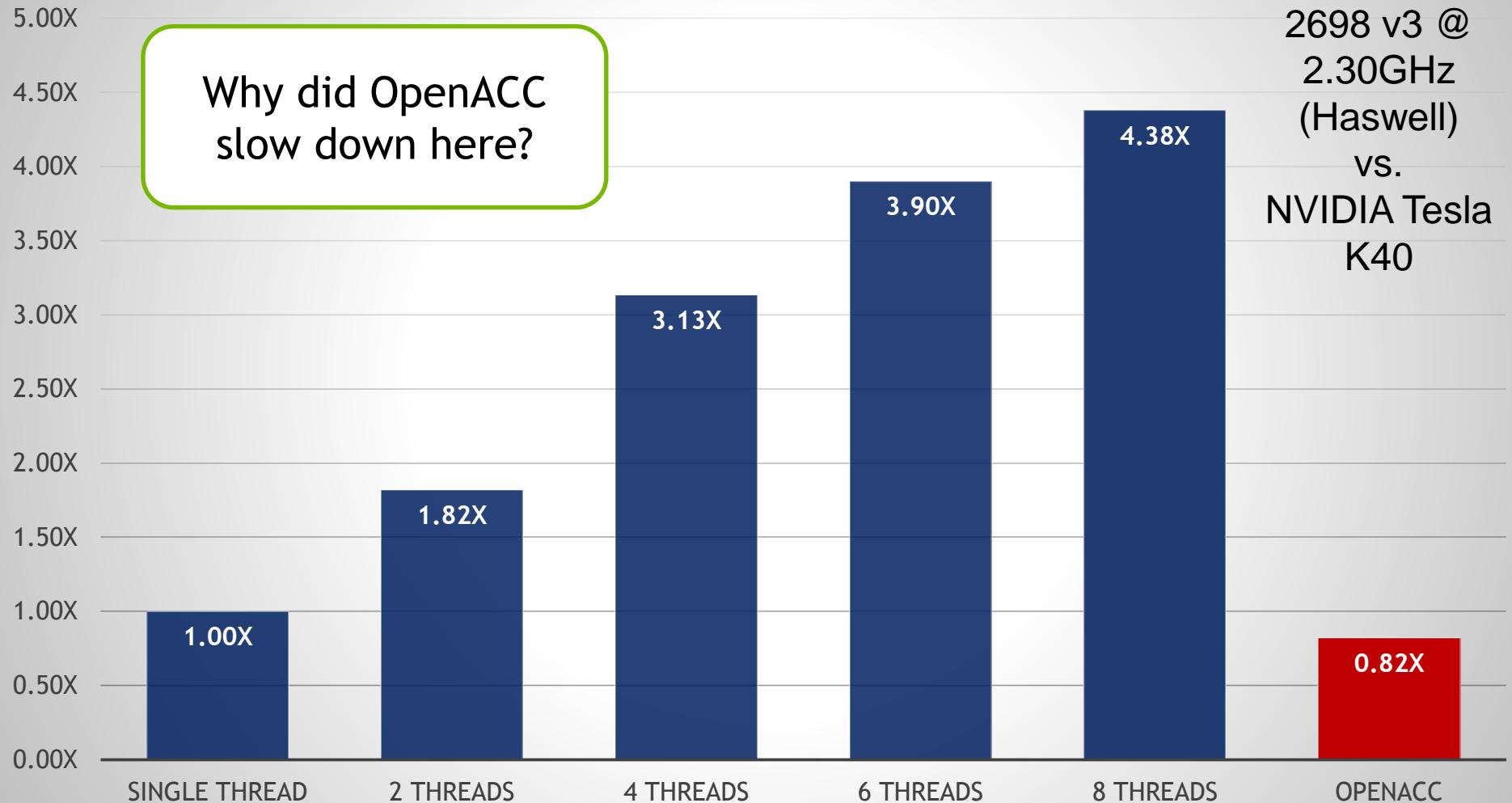
- Requires analysis by programmer to ensure safe parallelism
- Will parallelize what a compiler may miss
- Straightforward path from OpenMP

KERNELS

- Compiler performs parallel analysis and parallelizes what it believes safe
- Can cover larger area of code with single directive
- Gives compiler additional leeway to optimize.

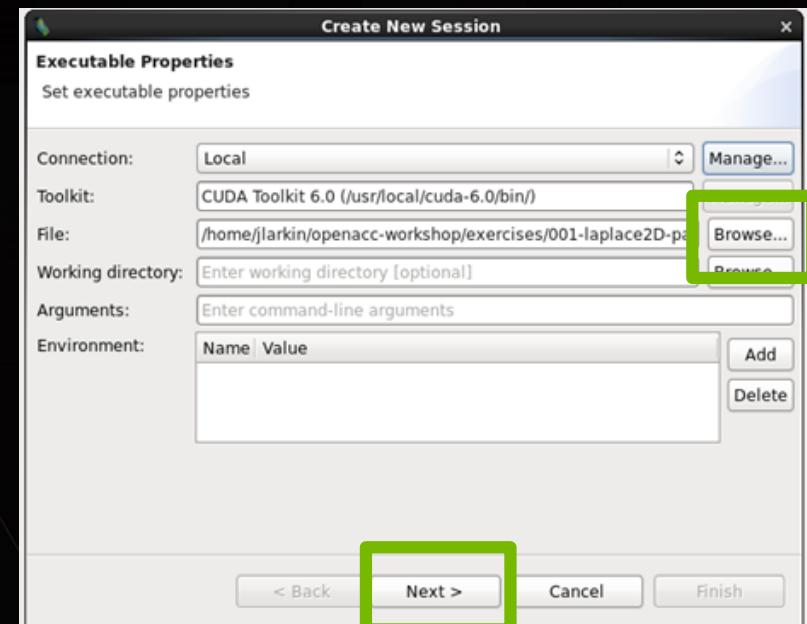
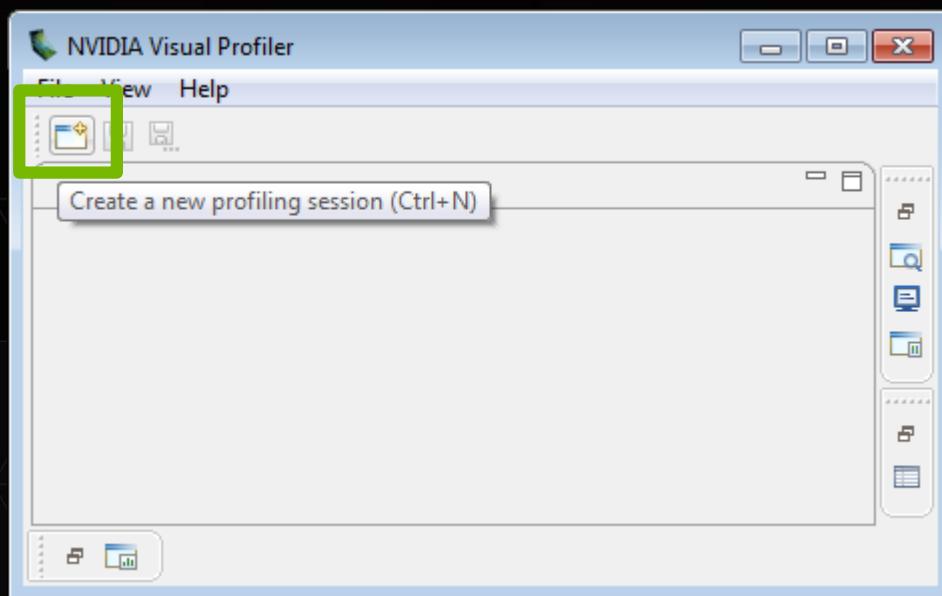
Both approaches are equally valid and can perform equally well.

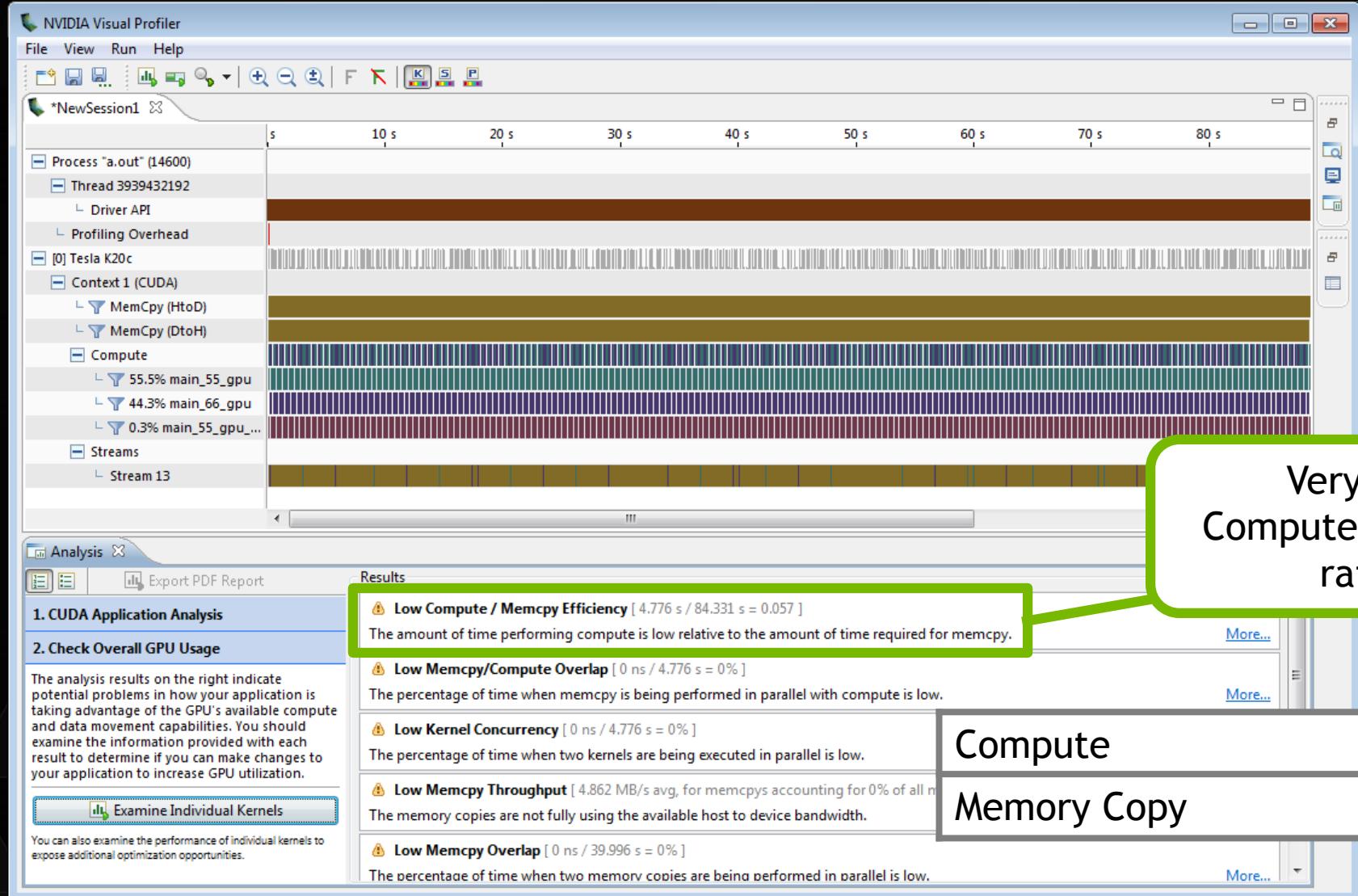
Speed-up (Higher is Better)



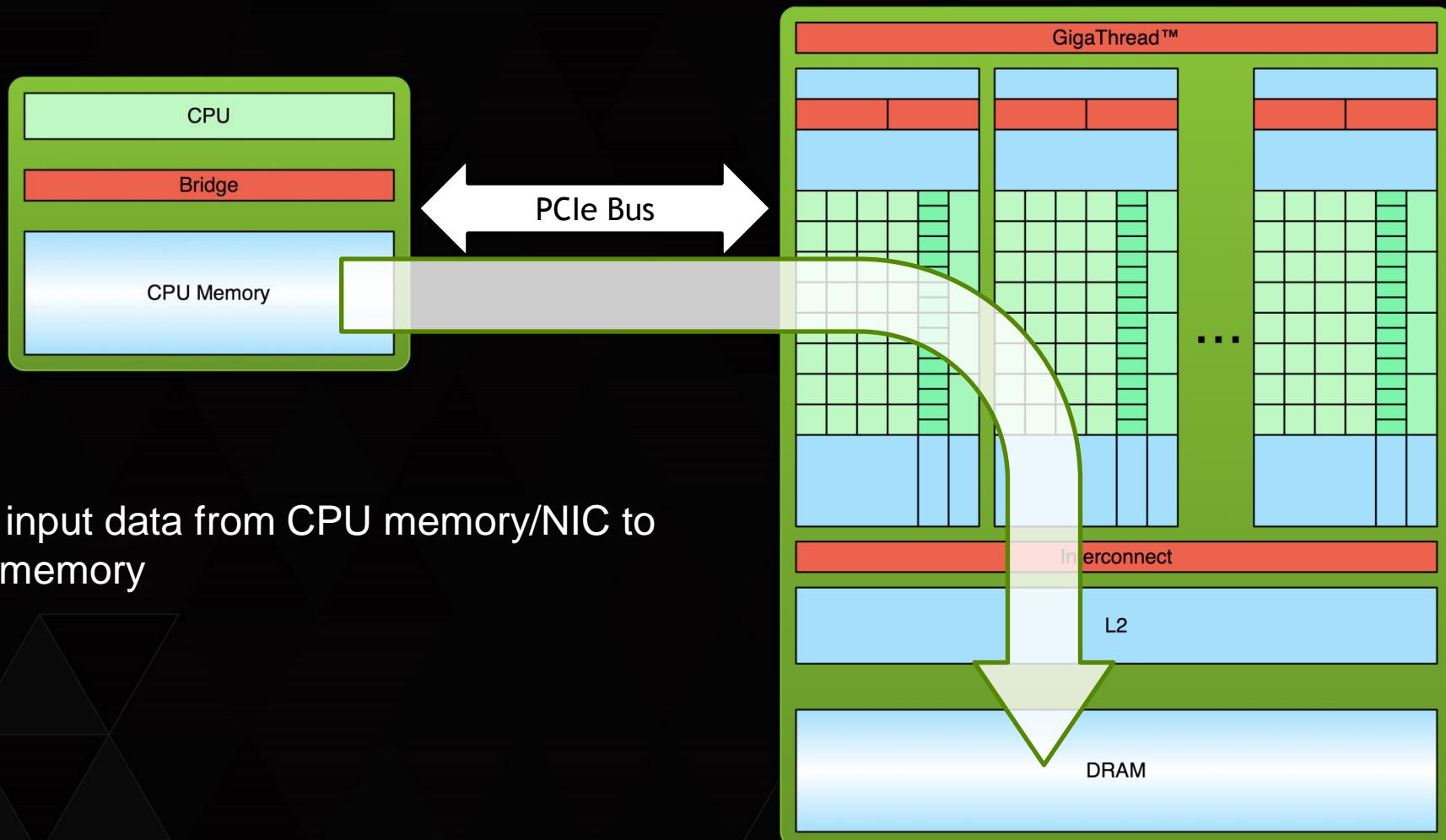
ANALYZING OPENACC PERFORMANCE

- ▶ Any tool that supports CUDA can likewise obtain performance information about OpenACC.
- ▶ Nvidia Visual Profiler (nvp) comes with the CUDA Toolkit, so it will be available on any machine with CUDA installed

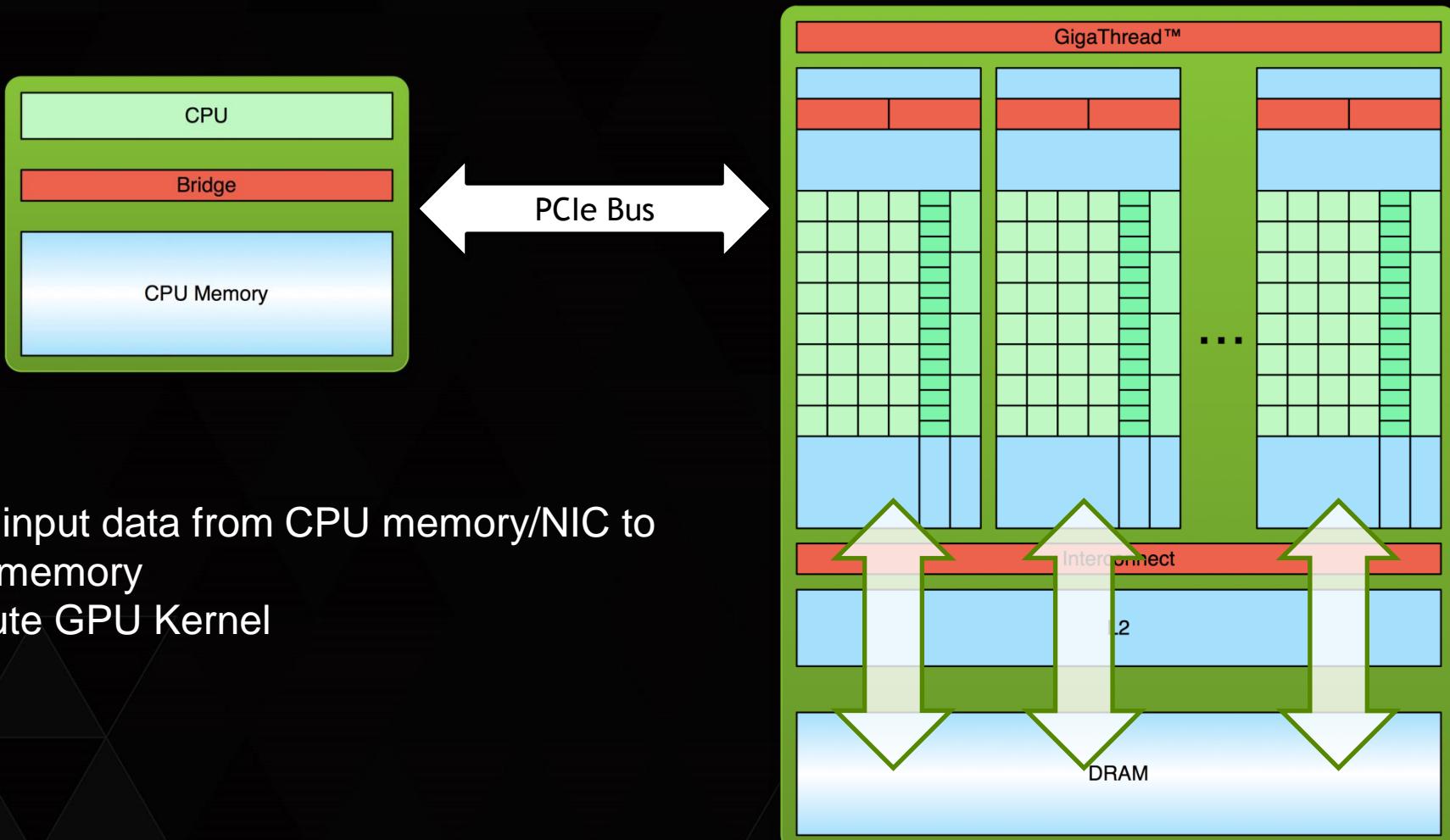




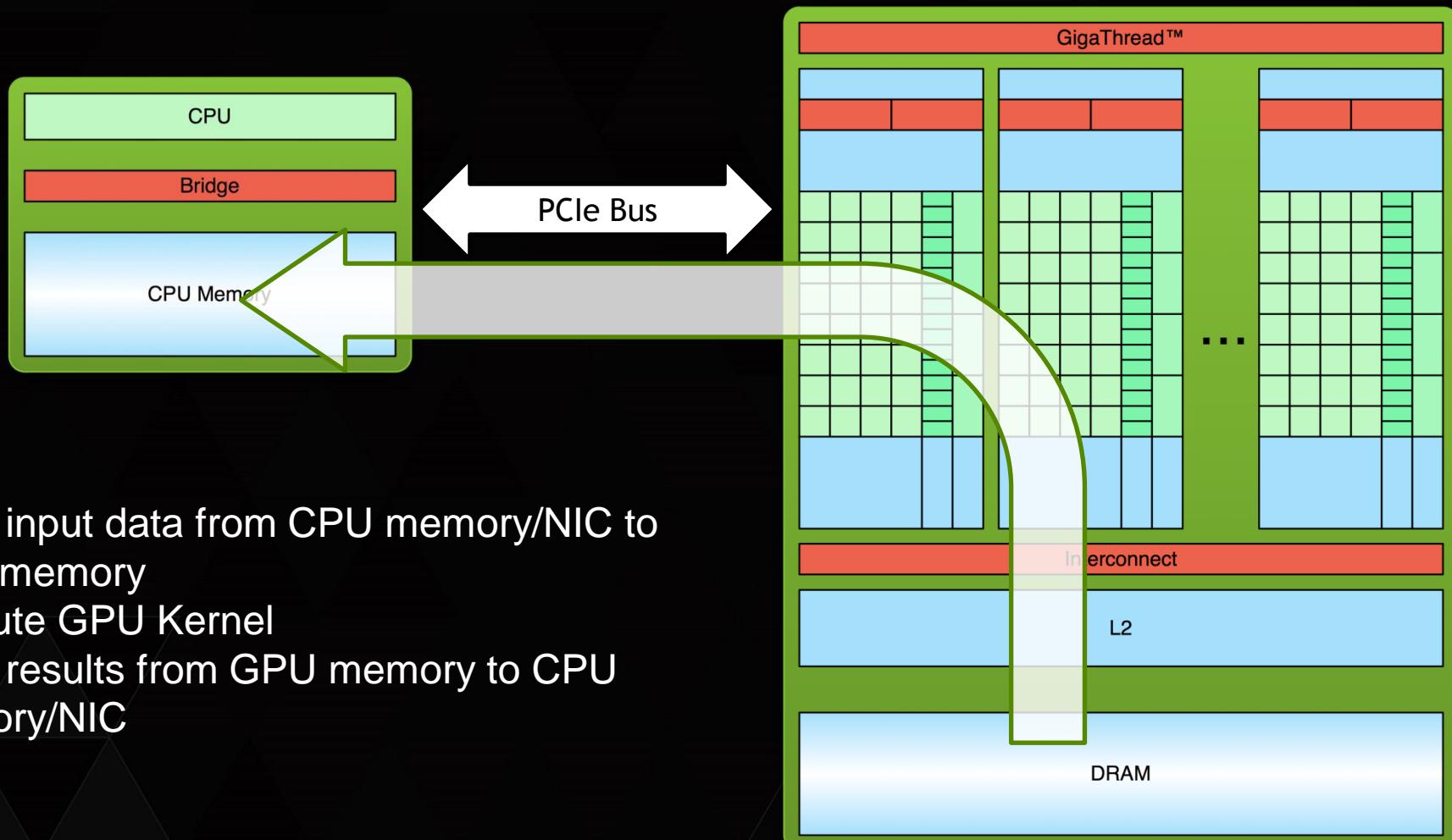
PROCESSING FLOW

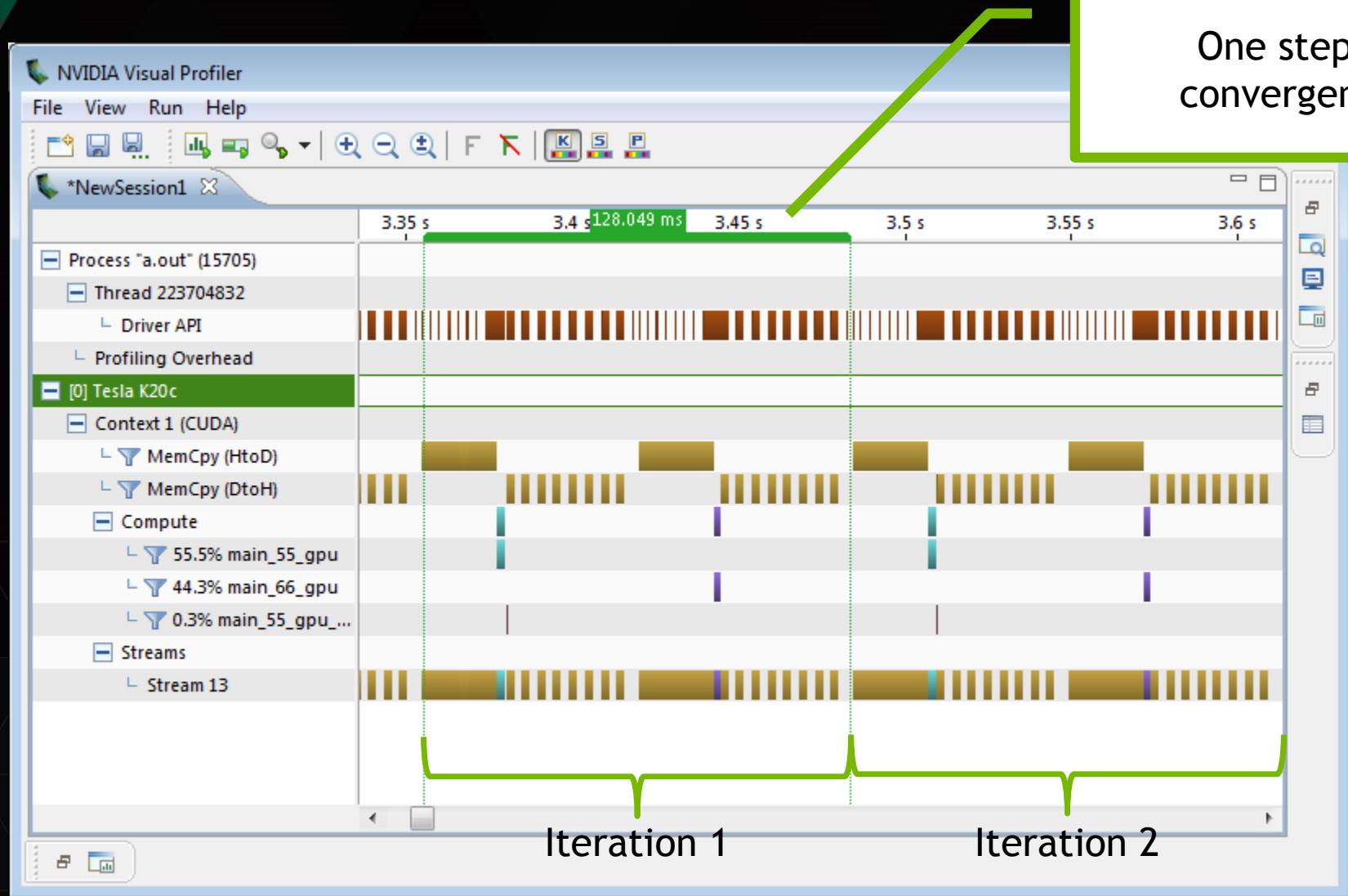


PROCESSING FLOW



PROCESSING FLOW





EXCESSIVE DATA TRANSFERS

```
while ( err > tol && iter < iter_max )  
{
```

```
    err=0.0;
```

A, Anew resident
on host

These copies
happen every
iteration of the
outer while
loop!*

A, Anew resident
on host

```
#pragma acc parallel loop reduction(max:err)
```

A, Anew resident on
accelerator

```
for( int j = 1; j < n-1; j++ ) {  
    for(int i = 1; i < m-1; i++) {  
        Anew[j][i] = 0.25 * (A[j][i+1] +  
                               A[j][i-1] + A[j-1][i] +  
                               A[j+1][i]);  
        err = max(err, abs(Anew[j][i] -  
                           A[j][i]));  
    }  
}
```

A, Anew resident on
accelerator

```
}
```

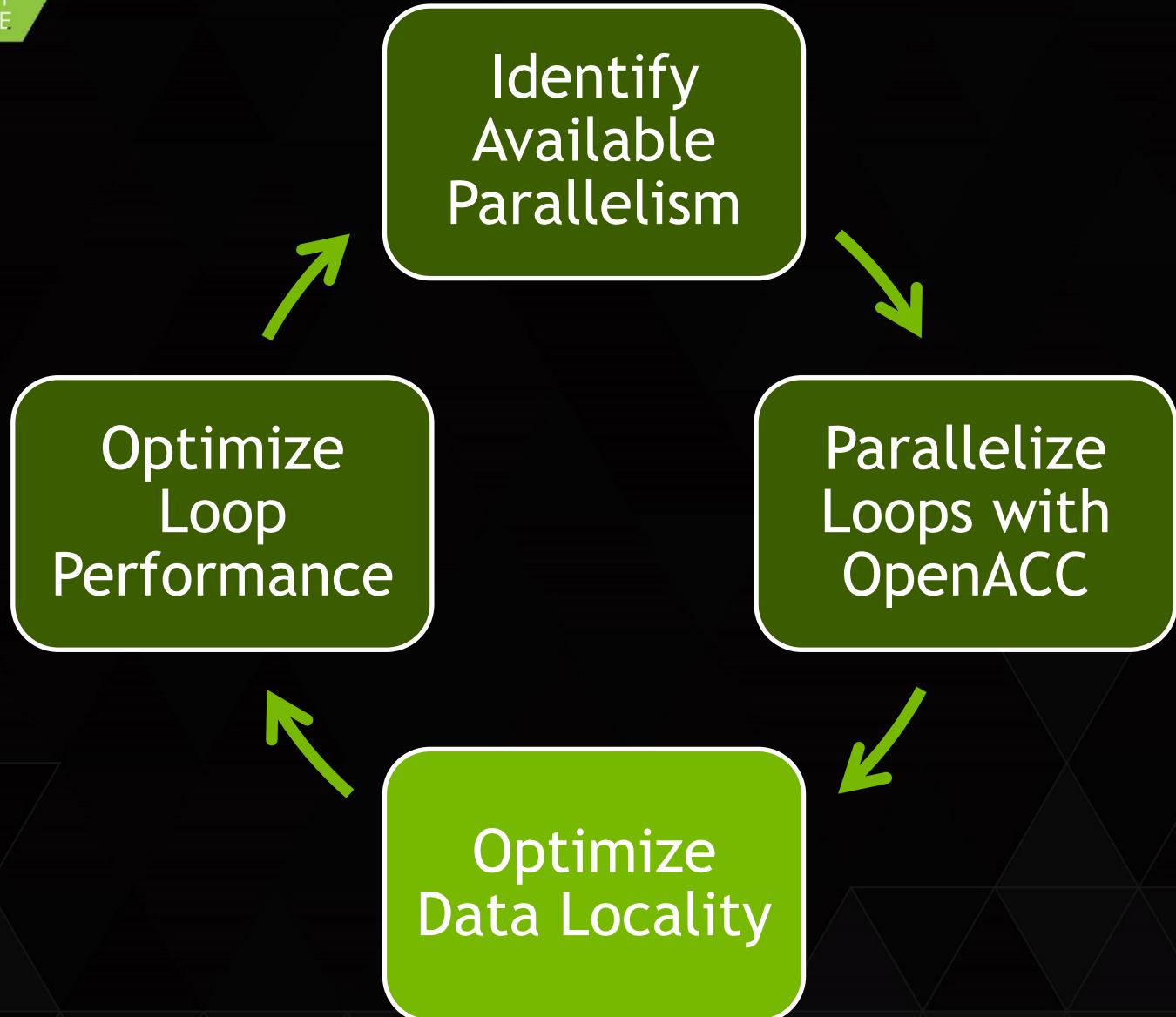
... And note that there are two `#pragma acc parallel`, so there are 4 copies per while loop iteration!

IDENTIFYING DATA LOCALITY

```
while ( err > tol && iter < iter_max ) {  
    err=0.0;  
  
#pragma acc parallel loop reduction(max:err)  
    for( int j = 1; j < n-1; j++) {  
        for(int i = 1; i < m-1; i++) {  
  
            Anew[j][i] = 0.25 * (A[j][i+1] + A[j][i-1] +  
                                  A[j-1][i] + A[j+1][i]);  
  
            err = max(err, abs(Anew[j][i] - A[j][i]));  
        }  
    }  
  
#pragma acc parallel loop  
    for( int j = 1; j < n-1; j++) {  
        for( int i = 1; i < m-1; i++ ) {  
            A[j][i] = Anew[j][i];  
        }  
    }  
  
    iter++;  
}
```

Does the CPU need the data between these loop nests?

Does the CPU need the data between iterations of the convergence loop?



DEFINING DATA REGIONS

- ▶ The **data** construct defines a region of code in which GPU arrays remain on the GPU and are shared among all kernels in that region.

```
#pragma acc data
{
#pragma acc parallel loop
...
#pragma acc parallel loop
...
}
```



Arrays used within the data region will remain on the GPU until the end of the data region.

DATA CLAUSES

- copy (list)** Allocates memory on GPU and copies data from host to GPU when entering region and copies data to the host when exiting region.
 - copyin (list)** Allocates memory on GPU and copies data from host to GPU when entering region.
 - copyout (list)** Allocates memory on GPU and copies data to the host when exiting region.
 - create (list)** Allocates memory on GPU but does not copy.
 - present (list)** Data is already present on GPU from another containing data region.
- and **present_or_copy[in|out]**, **present_or_create**, **deviceptr**.

The next OpenACC makes **present_or_*** the default behavior.

ARRAY SHAPING

- ▶ Compiler sometimes cannot determine size of arrays
 - ▶ Must specify explicitly using data clauses and array “shape”

C/C++

```
#pragma acc data copyin(a[0:size]),  
copyout(b[s/4:3*s/4])
```

Fortran

```
!$acc data copyin(a(1:end)),  
copyout(b(s/4:3*s/4))
```

- ▶ Note: data clauses can be used on **data**, **parallel**, or **kernels**

OPTIMIZE DATA LOCALITY

```
#pragma acc data copy(A) create(Anew)
while ( err > tol && iter < iter_max ) {
    err=0.0;

#pragma acc parallel loop reduction(max:err)
    for( int j = 1; j < n-1; j++ ) {
        for(int i = 1; i < m-1; i++) {

            Anew[j][i] = 0.25 * (A[j][i+1] + A[j][i-1] +
                                  A[j-1][i] + A[j+1][i]);

            err = max(err, abs(Anew[j][i] - A[j][i]));
        }
    }

#pragma acc parallel loop
    for( int j = 1; j < n-1; j++ ) {
        for( int i = 1; i < m-1; i++ ) {
            A[j][i] = Anew[j][i];
        }
    }

    iter++;
}
```

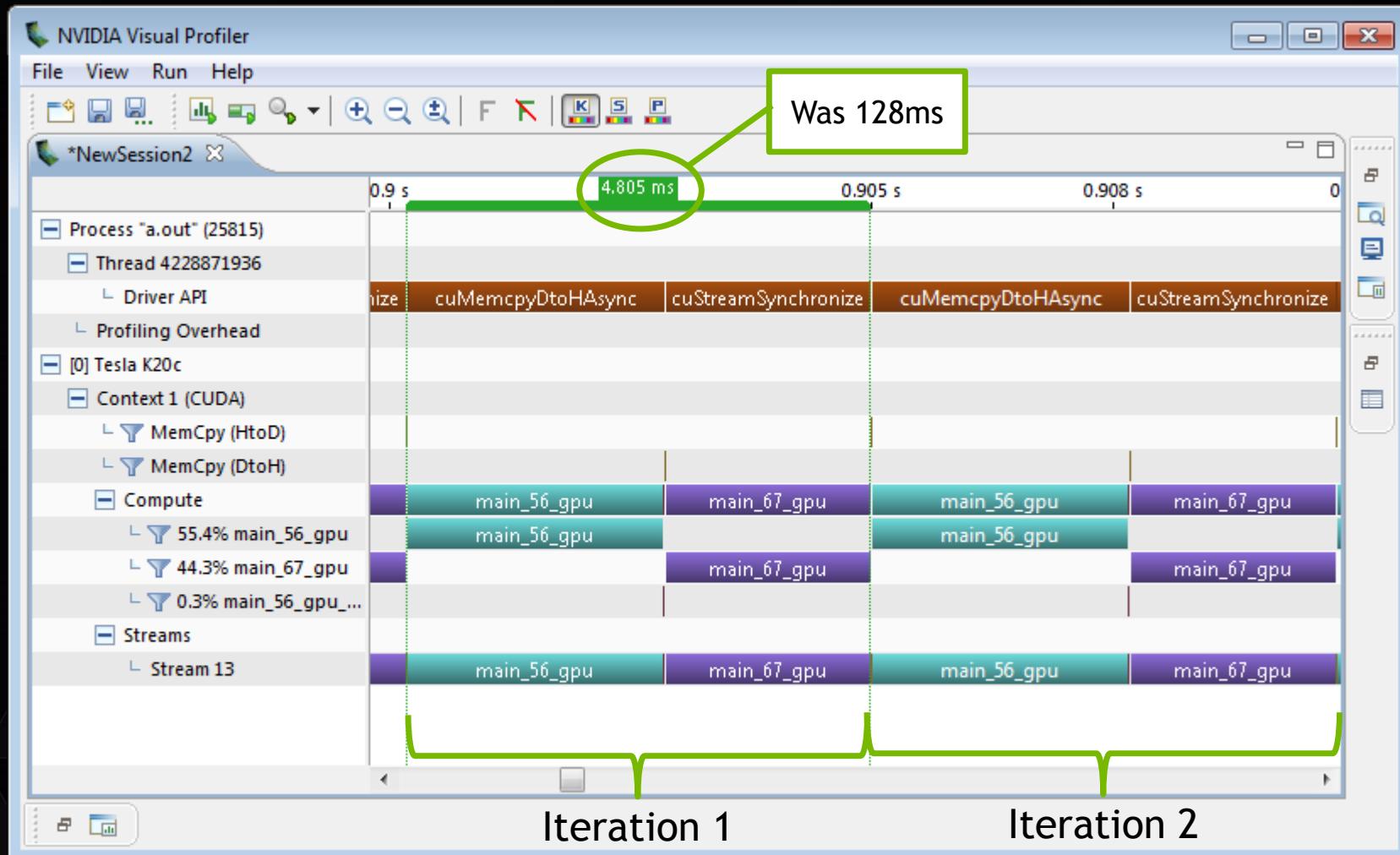
Copy A to/from the accelerator only when needed.

Create Anew as a device temporary.

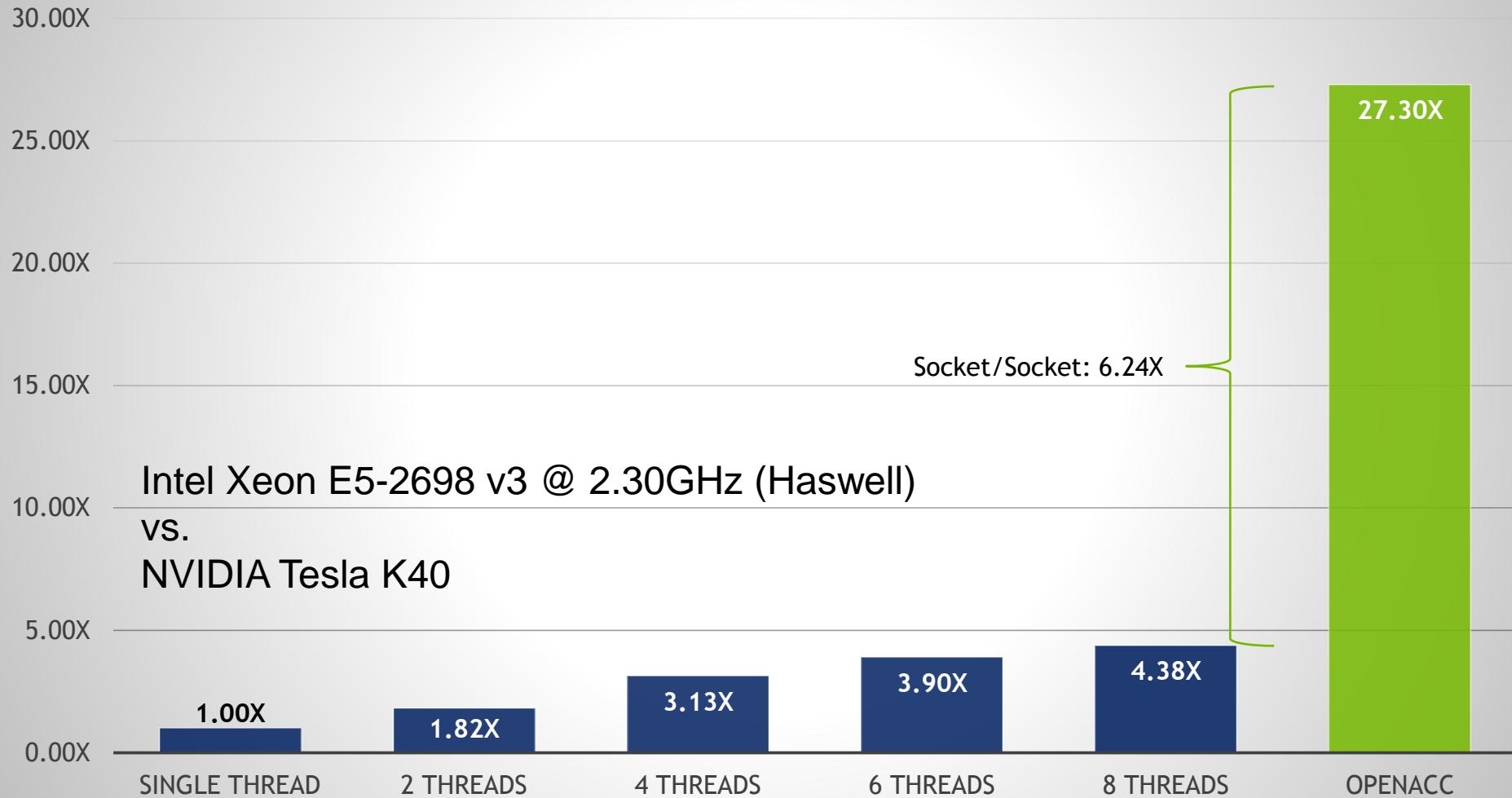
REBUILDING THE CODE

```
$ pgcc -fast -acc -ta=tesla -Minfo=all laplace2d.c
main:
  40, Loop not fused: function call before adjacent loop
    Generated vector sse code for the loop
  51, Generating copy(A[:, :])
    Generating create(Anew[:, :])
    Loop not vectorized/parallelized: potential early exits
  56, Accelerator kernel generated
    56, Max reduction generated for error
    57, #pragma acc loop gang /* blockIdx.x */
    59, #pragma acc loop vector(256) /* threadIdx.x */
  56, Generating Tesla code
  59, Loop is parallelizable
  67, Accelerator kernel generated
    68, #pragma acc loop gang /* blockIdx.x */
    70, #pragma acc loop vector(256) /* threadIdx.x */
  67, Generating Tesla code
  70, Loop is parallelizable
```

VISUAL PROFILER: DATA REGION



Speed-Up (Higher is Better)



OPENACC PRESENT CLAUSE

It's sometimes necessary for a data region to be in a different scope than the compute region.

When this occurs, the **present** clause can be used to tell the compiler data is already on the device.

Since the declaration of A is now in a higher scope, it's necessary to shape A in the present clause.

High-level data regions and the present clause are often critical to good performance.

```
function main(int argc, char **argv)
{
    #pragma acc data copy(A)
    {
        laplace2D(A,n,m);
    }
}
```

```
function laplace2D(double[N][M] A,n,m)
{
    #pragma acc data present(A[n][m]) create(Anew)
    while (err > tol && iter < iter_max) {
        err=0.0;
        ...
    }
}
```

UNSTRUCTURED DATA DIRECTIVES

Used to define data regions when scoping doesn't allow the use of normal data regions (e.g. The constructor/destructor of a class).

enter data Defines the start of an unstructured data lifetime

clauses: `copyin(list)`, `create(list)`

exit data Defines the end of an unstructured data lifetime

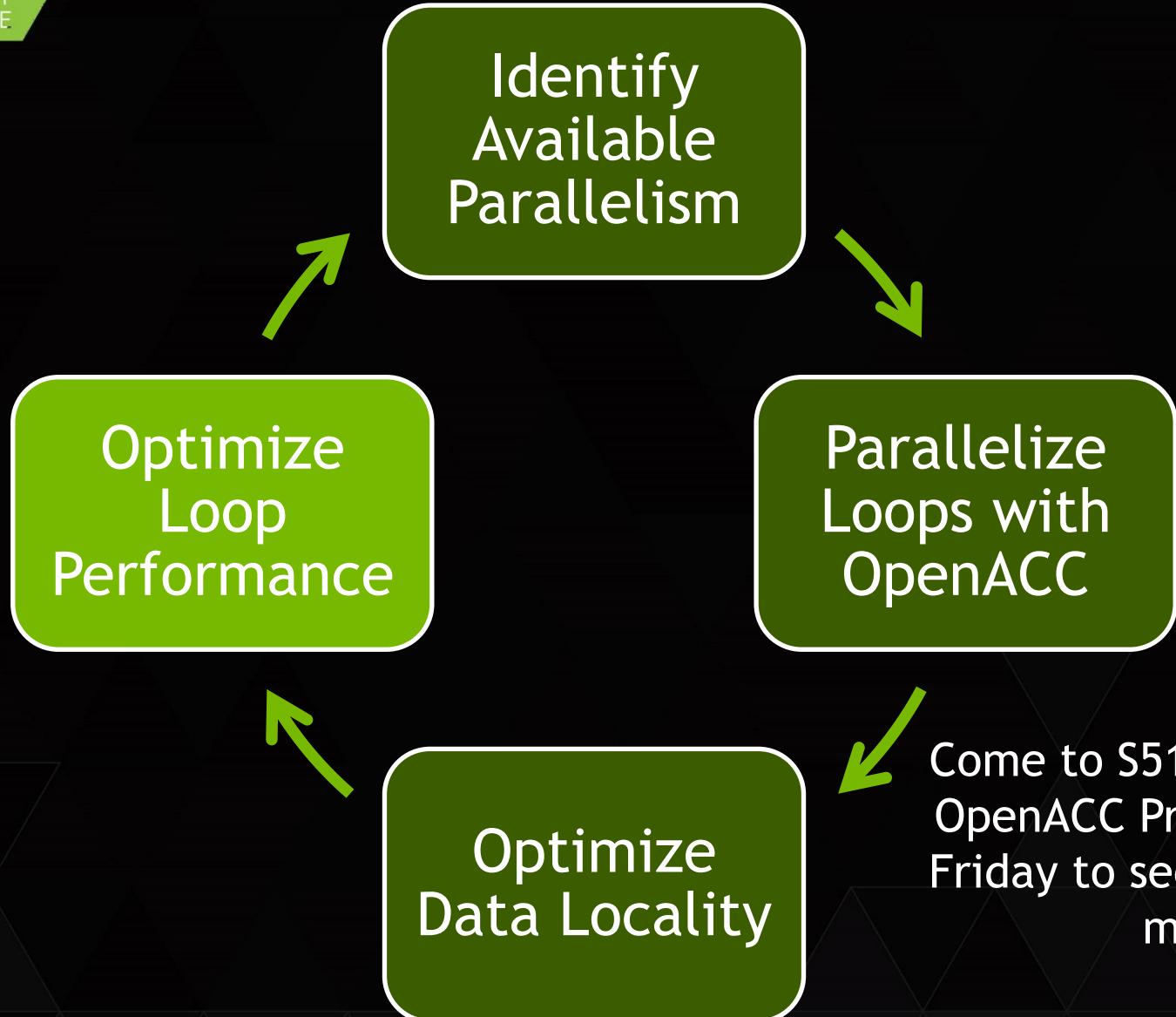
clauses: `copyout(list)`, `delete(list)`

```
#pragma acc enter data copyin(a)
...
#pragma acc exit data delete(a)
```

UNSTRUCTURED DATA REGIONS: C++ CLASSES

```
class Matrix {  
    Matrix(int n) {  
        len = n;  
        v = new double[len];  
        #pragma acc enter data create(v[0:len])  
    }  
    ~Matrix() {  
        #pragma acc exit data delete(v[0:len])  
        delete[] v;  
    }  
  
    private:  
    double* v;  
    int len;  
};
```

- ▶ Unstructured Data Regions enable OpenACC to be used in C++ classes
- ▶ Unstructured data regions can be used whenever data is allocated and initialized in a different scope than where it is freed.



MISC ADVICE

ALIASING CAN PREVENT PARALLELIZATION

23, Loop is parallelizable

Accelerator kernel generated

```
23, #pragma acc loop gang, vector(128) /* blockIdx.x threadIdx.x */
```

25, Complex loop carried dependence of 'b->' prevents parallelization

Loop carried dependence of 'a->' prevents parallelization

Loop carried backward dependence of 'a->' prevents vectorization

Accelerator scalar kernel generated

27, Complex loop carried dependence of 'a->' prevents parallelization

Loop carried dependence of 'b->' prevents parallelization

Loop carried backward dependence of 'b->' prevents vectorization

Accelerator scalar kernel generated

C99: RESTRICT KEYWORD

- ▶ Declaration of intent given by the programmer to the compiler

Applied to a pointer, e.g.

```
float *restrict ptr
```

Meaning: “for the lifetime of ptr, only it or a value directly derived from it (such as `ptr + 1`) will be used to access the object to which it points”*

- ▶ Parallelizing compilers often require `restrict` to determine independence
 - ▶ Otherwise the compiler can't parallelize loops that access `ptr`
 - ▶ Note: if programmer violates the declaration, behavior is undefined



~~float restrict *ptr~~

float *restrict ptr

<http://en.wikipedia.org/wiki/Restrict>

OPENACC INDEPENDENT CLAUSE

Specifies that loop iterations are data independent. This overrides any compiler dependency analysis. This is implied for *parallel loop*.

```
#pragma acc kernels
{
    #pragma acc loop independent
    for(int i=0; i<N; i++)
    {
        a[i] = 0.0;
        b[i] = 1.0;
        c[i] = 2.0;
    }
    #pragma acc loop independent
    for(int i=0; i<N; i++)
    {
        a(i) = b(i) + c(i)
    }
}
```



kernel 1

kernel 2

Informs the compiler that both loops are safe to parallelize so it will generate both kernels.

WRITE PARALLELIZABLE LOOPS

Use countable loops

C99: while->for

Fortran: while->do

Avoid pointer arithmetic

Write rectangular loops (compiler cannot parallelize triangular loops)

```
bool found=false;  
while(!found && i<N) {  
    if(a[i]==val) {  
        found=true  
        loc=i;  
    }  
    i++;  
}
```

```
bool found=false;  
for(int i=0;i<N;i++) {  
    if(a[i]==val) {  
        found=true  
        loc=i;  
    }  
}
```

```
for(int i=0;i<N;i++) {  
    for(int j=i;j<N;j++) {  
        sum+=A[i][j];  
    }  
}
```

```
for(int i=0;i<N;i++) {  
    for(int j=0;j<N;j++) {  
        if(j>=i)  
            sum+=A[i][j];  
    }  
}
```

OPENACC ROUTINE DIRECTIVE

The routine directive specifies that the compiler should generate a device copy of the function/subroutine in addition to the host copy and what type of parallelism the routine contains.

Clauses:

- ▶ **gang/worker/vector/seq**
 - ▶ Specifies the level of parallelism contained in the routine.
- ▶ **bind**
 - ▶ Specifies an optional name for the routine, also supplied at call-site
- ▶ **no_host**
 - ▶ The routine will only be used on the device
- ▶ **device_type**
 - ▶ Specialize this routine for a particular device type.

OPENACC DEBUGGING

- ▶ Most OpenACC directives accept an **if(condition)** clause

```
#pragma acc update self(A) if(debug)
#pragma acc parallel loop if(!debug)
[...]
#pragma acc update device(A) if(debug)
```

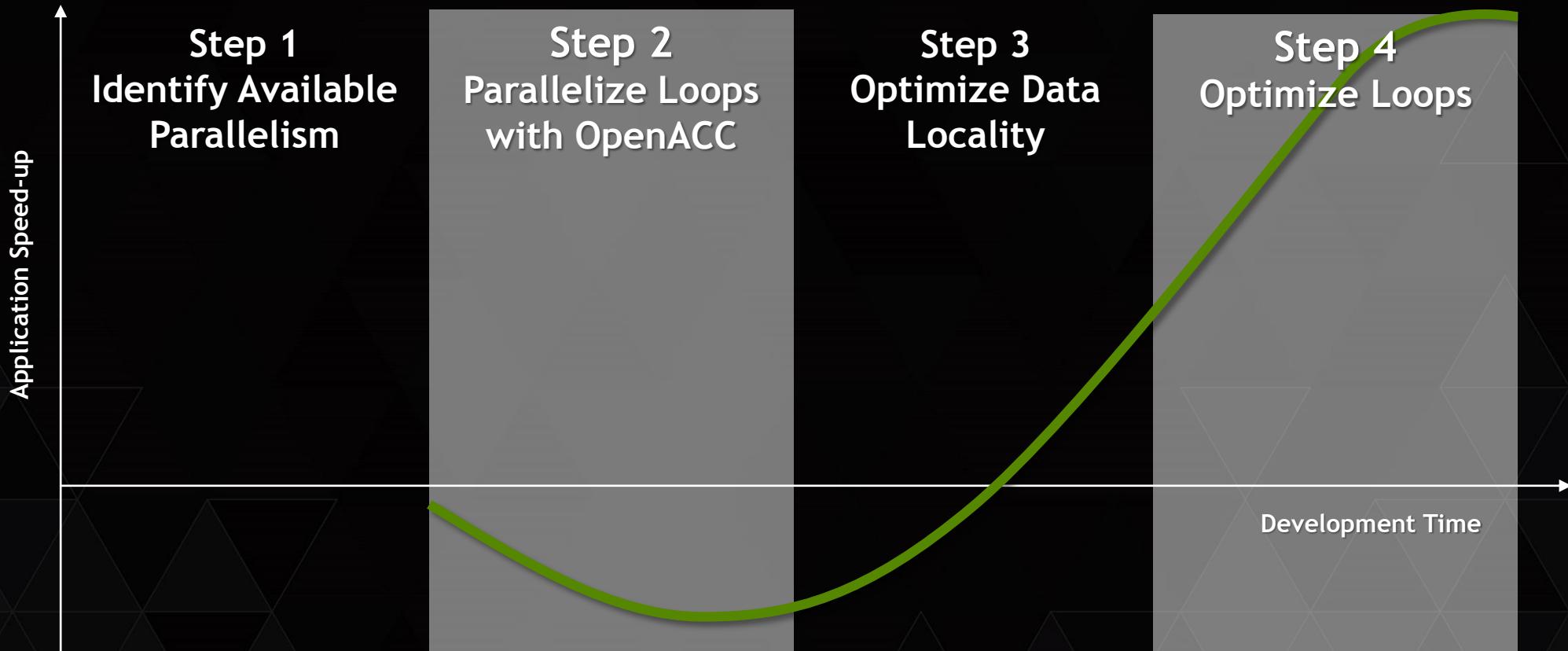
- ▶ Use **default(none)** to force explicit data directives

```
#pragma acc data copy(...) create(...) default(none)
```

NEXT STEPS

1. Identify Available Parallelism
 - ▶ What important parts of the code have available parallelism?
2. Parallelize Loops
 - ▶ Express as much parallelism as possible and ensure you still get correct results.
 - ▶ Because the compiler *must* be cautious about data movement, the code will generally slow down.
3. Optimize Data Locality
 - ▶ The programmer will *always* know better than the compiler what data movement is unnecessary.
4. Optimize Loop Performance
 - ▶ Don't try to optimize a kernel that runs in a few *us* or *ms* until you've eliminated the excess data motion that is taking *many seconds*.

TYPICAL PORTING EXPERIENCE WITH OPENACC DIRECTIVES



OPENACC AT GTC

S5192	Introduction to Compiler Directives w/ OpenACC	Wed 0900-1010	210H
S5388	OpenACC for Fortran Programmers	Wed 1400-1450	210H
S5139	Enabling OpenACC Performance Analysis	Wed 1500-1525	210H
S5515	Porting Apps to Titan: Results from the Hackathon	Wed 1600-1650	210H
S5233	GPU Acceleration Using OpenACC and C++ Classes	Thu 0900-0950	210D
S5382	OpenACC 2.5 and Beyond	Thu 1530-1555	220C
S5195	Advanced OpenACC Programming	Fri 0900-1020	210C
S5340	OpenACC and C++: An Application Perspective	Fri 1030-1055	210C
S5198	Panel on GPU Computing with OpenACC and OpenMP	Fri 1100-1150	210C

Plus many more sessions and OpenACC hang-outs!

NEXT STEPS

- ▶ Attend more OpenACC sessions at GTC.
- ▶ Try an OpenACC self-paced lab.
- ▶ Get a free trial of the PGI Compiler (www.pgroup.com)
- ▶ Please remember to fill out your surveys.

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